ECE: Information and Network Security Homework 5 Solutions Spring 2006

Chapter 15:

2. (a) There are four possible outcomes: HH, HT, TH, and TT. The probabilities are

$$p(HH) = p^2$$

 $p(HT) = p(1-p)$
 $p(TH) = p(1-p)$
 $p(TT) = (1-p)^2$

(b) The entropy is

$$-\left(p^2\log_2 p^2 + 2p(1-p)\log_2(p(1-p)) + (1-p)^2\log_2(1-p)^2\right).$$

By expanding and manipulating this can be expressed as $-2p \log_2 p - 2(1-p) \log_2(1-p)$. This could have been calculated easier using the fact that H(X,Y) = H(X) + H(Y) when X and Y are independent. Now observe that when X and Y are independent flips of the unfair coin, that $H(X,Y) = 2H(X) = -2p \log_2 p - 2(1-p) \log_2(1-p)$.

5. (a) $Y = 2^X$, so the possible outcomes for Y are 1/4, 1/2, 0, 2, and 4. The probabilities for Y are the same as the probabilities for X. Since the entropy only uses the probabilities, H(X) = H(Y).

(b) Observe that the possible outcomes for Y are 4, 1, and 0. In particular, two elements of X get mapped to Y = 4 and two elements get mapped to Y = 1. Since x^2 is not a one-to-one function, we have $H(Y) \leq H(X)$.

9. (a) The plaintext entropy is

$$H(P) = -\left(\frac{1}{3}\log_2\frac{1}{3} + \frac{2}{3}\log_2\frac{2}{3}\right).$$

(b) Observe that this system matches up with the one-time pad, and hence H(P|C) = H(P).

10. (a) H(P, K) = H(C, P, K) since knowledge of the plaintext and the key determine the ciphertext. H(P, K) = H(P) + H(K) since the keys are chosen independently of the plaintext in a cryptosystem.

(b) H(C, P) = H(C) + H(P|C) by the Chain Rule. Since we have perfect secrecy, H(P|C) = H(P), and thus H(C, P) = H(C) + H(P). For the last part, refer to the solution to problem 14.11 to get H(C|P) = H(K) - H(K|C, P). We must show that in a system with perfect secrecy, that H(C|P) =H(C). To see this, observe that H(C, P) = H(C) + H(P|C) = H(P) + H(C|P). Since H(P) = H(P|C), we have H(C) = H(C|P).

(c) Use the last part of (b) and observe that the stated condition implies that H(K|C, P) = 0.