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## Chapter 9

# Differential Entropy

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1. *Differential entropy.* Evaluate the differential entropy  $h(X) = -\int f \ln f$  for the following:

- (a) The exponential density,  $f(x) = \lambda e^{-\lambda x}$ ,  $x \geq 0$ .
- (b) The Laplace density,  $f(x) = \frac{1}{2} \lambda e^{-\lambda |x|}$ .
- (c) The sum of  $X_1$  and  $X_2$ , where  $X_1$  and  $X_2$  are independent normal random variables with means  $\mu_i$  and variances  $\sigma_i^2$ ,  $i = 1, 2$ .

**Solution:** *Differential Entropy.*

- (a) Exponential distribution.

$$h(f) = -\int_0^{\infty} \lambda e^{-\lambda x} [\ln \lambda - \lambda x] dx \quad (9.1)$$

$$= -\ln \lambda + 1 \text{ nats.} \quad (9.2)$$

$$= \log \frac{e}{\lambda} \text{ bits.} \quad (9.3)$$

- (b) Laplace density.

$$h(f) = -\int_{-\infty}^{\infty} \frac{1}{2} \lambda e^{-\lambda |x|} \left[ \ln \frac{1}{2} + \ln \lambda - \lambda |x| \right] dx \quad (9.4)$$

$$= -\ln \frac{1}{2} - \ln \lambda + 1 \quad (9.5)$$

$$= \ln \frac{2e}{\lambda} \text{ nats.} \quad (9.6)$$

$$= \log \frac{2e}{\lambda} \text{ bits.} \quad (9.7)$$

(c) Sum of two normal distributions.

The sum of two normal random variables is also normal, so applying the result derived the class for the normal distribution, since  $X_1 + X_2 \sim \mathcal{N}(\mu_1 + \mu_2, \sigma_1^2 + \sigma_2^2)$ ,

$$h(f) = \frac{1}{2} \log 2\pi e(\sigma_1^2 + \sigma_2^2) \text{ bits.} \quad (9.8)$$

2. *Concavity of determinants.* Let  $K_1$  and  $K_2$  be two symmetric nonnegative definite  $n \times n$  matrices. Prove the result of Ky Fan[4]:

$$|\lambda K_1 + \bar{\lambda} K_2| \geq |K_1|^\lambda |K_2|^{\bar{\lambda}}, \quad \text{for } 0 \leq \lambda \leq 1, \bar{\lambda} = 1 - \lambda,$$

where  $|K|$  denotes the determinant of  $K$ .

*Hint:* Let  $\mathbf{Z} = \mathbf{X}_\theta$ , where  $\mathbf{X}_1 \sim N(0, K_1)$ ,  $\mathbf{X}_2 \sim N(0, K_2)$  and  $\theta = \text{Bernoulli}(\lambda)$ . Then use  $H(\mathbf{Z} | \theta) \leq H(\mathbf{Z})$ .

*Solution: Concavity of Determinants.* Let  $X_1$  and  $X_2$  be normally distributed  $n$ -vectors,  $\mathbf{X}_i \sim \phi_{K_i}(\mathbf{x})$ ,  $i = 1, 2$ . Let the random variable  $\theta$  have distribution  $\Pr\{\theta = 1\} = \lambda$ ,  $\Pr\{\theta = 2\} = 1 - \lambda$ ,  $0 \leq \lambda \leq 1$ . Let  $\theta$ ,  $\mathbf{X}_1$ , and  $\mathbf{X}_2$  be independent and let  $\mathbf{Z} = \mathbf{X}_\theta$ . Then  $\mathbf{Z}$  has covariance  $K_Z = \lambda K_1 + (1 - \lambda) K_2$ . However,  $\mathbf{Z}$  will not be multivariate normal. However, since a normal distribution maximizes the entropy for a given variance, we have

$$\frac{1}{2} \ln(2\pi e)^n |\lambda K_1 + (1 - \lambda) K_2| \geq h(\mathbf{Z}) \geq h(\mathbf{Z} | \theta) = \lambda \frac{1}{2} \ln(2\pi e)^n |K_1| + (1 - \lambda) \frac{1}{2} \ln(2\pi e)^n |K_2|. \quad (9.9)$$

Thus

$$|\lambda K_1 + (1 - \lambda) K_2| \geq |K_1|^\lambda |K_2|^{1 - \lambda}, \quad (9.10)$$

as desired.

3. *Mutual information for correlated normals.* Find the mutual information  $I(X; Y)$ , where

$$\begin{pmatrix} X \\ Y \end{pmatrix} \sim N_2 \left( 0, \begin{bmatrix} \sigma^2 & \rho\sigma^2 \\ \rho\sigma^2 & \sigma^2 \end{bmatrix} \right).$$

Evaluate  $I(X; Y)$  for  $\rho = 1$ ,  $\rho = 0$ , and  $\rho = -1$ , and comment.

*Solution: Mutual information for correlated normals.*

$$\begin{bmatrix} X \\ Y \end{bmatrix} \sim \mathcal{N}_2 \left( 0, \begin{bmatrix} \sigma^2 & \rho\sigma^2 \\ \rho\sigma^2 & \sigma^2 \end{bmatrix} \right) \quad (9.11)$$

Using the expression for the entropy of a multivariate normal derived in class

$$h(X, Y) = \frac{1}{2} \log(2\pi e)^2 |K| = \frac{1}{2} \log(2\pi e)^2 \sigma^4 (1 - \rho^2). \quad (9.12)$$

Since  $X$  and  $Y$  are individually normal with variance  $\sigma^2$ ,

$$h(X) = h(Y) = \frac{1}{2} \log 2\pi e \sigma^2. \quad (9.13)$$

Hence

$$I(X;Y) = h(X) + h(Y) - h(X,Y) = -\frac{1}{2} \log(1 - \rho^2). \quad (9.14)$$

- (a)  $\rho = 1$ . In this case,  $X = Y$ , and knowing  $X$  implies perfect knowledge about  $Y$ . Hence the mutual information is infinite, which agrees with the formula.
- (b)  $\rho = 0$ . In this case,  $X$  and  $Y$  are independent, and hence  $I(X;Y) = 0$ , which agrees with the formula.
- (c)  $\rho = -1$ . In this case,  $X = -Y$ , and again the mutual information is infinite as in the case when  $\rho = 1$ .
4. *Uniformly distributed noise.* Let the input random variable  $X$  to a channel be uniformly distributed over the interval  $-1/2 \leq x \leq +1/2$ . Let the output of the channel be  $Y = X + Z$ , where the noise random variable is uniformly distributed over the interval  $-a/2 \leq z \leq +a/2$ .
- (a) Find  $I(X;Y)$  as a function of  $a$ .
- (b) For  $a = 1$  find the capacity of the channel when the input  $X$  is peak-limited; that is, the range of  $X$  is limited to  $-1/2 \leq x \leq +1/2$ . What probability distribution on  $X$  maximizes the mutual information  $I(X;Y)$ ?
- (c) (Optional) Find the capacity of the channel for all values of  $a$ , again assuming that the range of  $X$  is limited to  $-1/2 \leq x \leq +1/2$ .

**Solution:** *Uniformly distributed noise.* The probability density function for  $Y = X + Z$  is the convolution of the densities of  $X$  and  $Z$ . Since both  $X$  and  $Z$  have rectangular densities, the density of  $Y$  is a trapezoid. For  $a < 1$  the density for  $Y$  is

$$p_Y(y) = \begin{cases} (1/2a)(y + (1+a)/2) & -(1+a)/2 \leq y \leq -(1-a)/2 \\ 1 & -(1-a)/2 \leq y \leq +(1-a)/2 \\ (1/2a)(-y - (1+a)/2) & +(1-a)/2 \leq y \leq +(1+a)/2 \end{cases}$$

and for  $a > 1$  the density for  $Y$  is

$$p_Y(y) = \begin{cases} y + (a+1)/2 & -(a+1)/2 \leq y \leq -(a-1)/2 \\ 1/a & -(a-1)/2 \leq y \leq +(a-1)/2 \\ -y - (a+1)/2 & +(a-1)/2 \leq y \leq +(a+1)/2 \end{cases}$$

(When  $a = 1$ , the density of  $Y$  is triangular over the interval  $[-1, +1]$ .)

- (a) We use the identity  $I(X;Y) = h(Y) - h(Y|X)$ . It is easy to compute  $h(Y)$  directly, but it is even easier to use the grouping property of entropy. First suppose that  $a < 1$ . With probability  $1 - a$ , the output  $Y$  is conditionally uniformly distributed in the interval  $[-(1-a)/2, +(1-a)/2]$ ; whereas with probability  $a$ ,  $Y$  has a split triangular density where the base of the triangle has width  $a$ . As shown in examples in class,

$$\begin{aligned} h(Y) &= H(a) + (1-a) \ln(1-a) + a \left( \frac{1}{2} + \ln a \right) \\ &= -a \ln a - (1-a) \ln(1-a) + (1-a) \ln(1-a) + \frac{a}{2} + a \ln a = \frac{a}{2} \text{ nats.} \end{aligned}$$

If  $a > 1$  the trapezoidal density of  $Y$  can be scaled by a factor  $a$ , which yields  $h(Y) = \ln a + 1/2a$ . Given any value of  $x$ , the output  $Y$  is conditionally uniformly distributed over an interval of length  $a$ , so the conditional differential entropy in nats is  $h(Y|X) = h(Z) = \ln a$  for all  $a > 0$ . Therefore the mutual information in nats is

$$I(X;Y) = \begin{cases} a/2 - \ln a & \text{if } a \leq 1 \\ 1/2a & \text{if } a \geq 1. \end{cases}$$

As expected,  $I(X;Y) \rightarrow \infty$  as  $a \rightarrow 0$  and  $I(X;Y) \rightarrow 0$  as  $a \rightarrow \infty$ .

- (b) As usual with additive noise, we can express  $I(X;Y)$  in terms of  $h(Y)$  and  $h(Z)$ :

$$I(X;Y) = h(Y) - h(Y|X) = h(Y) - h(Z).$$

Since both  $X$  and  $Z$  are limited to the interval  $[-1/2, +1/2]$ , their sum  $Y$  is limited to the interval  $[-1, +1]$ . The differential entropy of  $Y$  is at most that of a random variable uniformly distributed on that interval; that is,  $h(Y) \leq 1$ . This maximum entropy can be achieved if the input  $X$  takes on its extreme values  $x = \pm 1$  each with probability  $1/2$ . In this case,  $I(X;Y) = h(Y) - h(Z) = 1 - 0 = 1$ . Decoding for this channel is quite simple:

$$\hat{X} = \begin{cases} -1/2 & \text{if } y < 0 \\ +1/2 & \text{if } y \geq 0. \end{cases}$$

This coding scheme transmits one bit per channel use with zero error probability. (Only a received value  $y = 0$  is ambiguous, and this occurs with probability 0.)

- (c) When  $a$  is of the form  $1/m$  for  $m = 2, 3, \dots$ , we can achieve the maximum possible value  $I(X;Y) = \log m$  when  $X$  is uniformly distributed over the discrete points  $\{-1, -1 + 2/(m-1), \dots, +1 - 2/(m-1), +1\}$ . In this case  $Y$  has a uniform probability density on the interval  $[-1 - 1/(m-1), +1 + 1/(m-1)]$ . Other values of  $a$  are left as an exercise.
5. *Quantized random variables.* Roughly how many bits are required on the average to describe to 3 digit accuracy the decay time (in years) of a radium atom if the half-life of radium is 80 years? Note that half-life is the median of the distribution.

**Solution:** *Quantized random variables.* The differential entropy of an exponentially distributed random variable with mean  $1/\lambda$  is  $\log \frac{e}{\lambda}$  bits. If the median is 80 years, then

$$\int_0^{80} \lambda e^{-\lambda x} dx = \frac{1}{2} \quad (9.15)$$

or

$$\lambda = \frac{\ln 2}{80} = 0.00866 \quad (9.16)$$

and the differential entropy is  $\log e/\lambda$ . To represent the random variable to 3 digits  $\approx 10$  bits accuracy would need  $\log e/\lambda + 10$  bits = 18.3 bits.

6. *Scaling.* Let  $h(\mathbf{X}) = -\int f(\mathbf{x}) \log f(\mathbf{x}) d\mathbf{x}$ . Show  $h(A\mathbf{X}) = \log |\det(A)| + h(\mathbf{X})$ .

**Solution:** *Scaling.* Let  $\mathbf{Y} = A\mathbf{X}$ . Then the density of  $\mathbf{Y}$  is

$$g(\mathbf{y}) = \frac{1}{|A|} f(A^{-1}\mathbf{y}). \quad (9.17)$$

Hence

$$h(A\mathbf{X}) = -\int g(\mathbf{y}) \ln g(\mathbf{y}) d\mathbf{y} \quad (9.18)$$

$$= -\int \frac{1}{|A|} f(A^{-1}\mathbf{y}) [\ln f(A^{-1}\mathbf{y}) - \log |A|] d\mathbf{y} \quad (9.19)$$

$$= -\int \frac{1}{|A|} f(\mathbf{x}) [\ln f(\mathbf{x}) - \log |A|] |A| d\mathbf{x} \quad (9.20)$$

$$= h(\mathbf{X}) + \log |A|. \quad (9.21)$$