

332:559

**Spectrum Management and Wireless Systems**  
Take Home Examination

**Fall 2003**

*Each problem PART is worth the same number of points and the whole exam is half of the grade – the other half will be determined from your conference paper proposals.*

1. **Eigenfunctions and Signal Space** Suppose a waveform  $\phi_k(t)$  is applied to linear time-invariant filter with impulse response  $h(t)$  and output  $r(t)$ .

- (a) If  $\phi_k(t)$  is zero outside  $t \in [0, T]$  and we only consider the output on the same interval, find an expression which can be solved to find the set of all  $\phi_k(t)$  such that

$$r(t) = (\phi_k * h)(t) = \lambda_k \phi_k(t)$$

where  $\lambda_k \in \mathcal{C}$ .

Such  $\phi_k(t)$  are called eigenfunctions of the system.

- (b) For  $h(t) = e^{-t}$  provide numerical examples of a few  $\phi_k(t)$  and the associated  $\lambda_k$  for  $T = 1$ . Plot the two with the largest eigenvalues and the two with the smallest eigenvalues. You may numerically calculate any integrals.
- (c) Suppose  $\phi_k(t)$  and  $r(t)$  are now defined  $\forall t$ . Show (or disprove) that  $e^{st}$  is an eigenfunction of any  $h(t)$  where  $s \in \mathcal{C}$ .
- (d) Now suppose we require the  $\phi_k(t)$  to be orthonormal. Then we cannot guarantee that  $\phi_k(t)$  is an eigenfunction of the system  $h(t)$ . However, if we define  $\psi_k(t)$  as the output associated with application of  $\phi_k(t)$  to  $h(t)$ , then we can require

$$\langle \phi_k(t), \phi_\ell(t) \rangle = \delta_{k\ell}$$

and

$$\langle \psi_k(t), \psi_\ell(t) \rangle = \lambda_k \delta_{k\ell}$$

Derive an expression from which the  $\phi_k(t)$  and thence the  $\psi_k(t)$  can be found. Assume the  $\phi_k(t)$  and  $\psi_k(t)$  are defined only on  $(0, T)$ .

- (e) Let  $h(t) = e^{-t}u(t)$ . Derive an expression for the  $\phi_k(t)$ , find a few of them numerically. Plot the two with the largest eigenvalues and the two with the smallest eigenvalues. Assume  $T = 1$ .
- (f) Let  $h(t) = \frac{\sin 2\pi t}{\pi t}$  and repeat the previous part for  $T = 1$  and then  $T = 100$ .
- (g) Let  $\phi_k(t)$  and  $\psi_k(t)$  now be defined  $\forall t$ . Analytically find appropriate  $\phi_k(t)$  for arbitrary  $h(t)$ .

2. **Waterfilling Math** You are given a power budget  $P$  which you can split over  $N$  Gaussian channels. Let the noise level in each channel be  $\sigma_i$  and the power used on that channel  $p_i$ . The capacity of each channel is

$$C_i = \frac{1}{2} \log \left( 1 + \frac{p_i}{\sigma_i} \right)$$

and the total capacity using all channels is

$$C = \sum_{i=1}^N C_i = \frac{1}{2} \sum_{i=1}^N \log \left( 1 + \frac{p_i}{\sigma_i} \right)$$

- (a) Suppose you are only allowed to place power into a single dimension, but you can choose which one. Show that any choice of dimension  $k$  must satisfy  $\sigma_k \leq \sigma_\ell$ ,  $\ell = 1, 2, \dots, N$  to attain maximum capacity  $C_k$ .
- (b) Prove using Lagrange Multipliers that

$$p_i = [c - \sigma_i]^+$$

where  $c$  is a constant chosen to satisfy  $\sum_i p_i = P$ .

Derive a condition on  $P$  so that no  $p_i$  is zero.

You can find writeups on Lagrange Multipliers in most optimization texts and some applied math texts. My favorite is *Advanced Calculus for Applications* (Hildebrand).

- (c) Let  $d_i = p_i + \sigma_i$ . Show using Jensen's inequality that to maximize capacity the  $d_i$  should be a constant  $c$ , or if not,  $\sigma_i$ . What is the constant  $c$ ? Is this the same result as the previous part?

HINT: This is a little of a mind-bender. For those of you who've done a formal optimization course, this is an example of optimization through a "slack variable" (I think). Once you get the basic idea, it might help when finishing up to recognize that  $\log 1 = 0$  and  $\log d_i \geq 0$  by the definition of  $d_i$  and capacity.

3. **Single Base Interference Avoidance** In class and in the assigned Interference Avoidance paper (see web page) we learned the following facts:

$$\mathbf{S} = \begin{bmatrix} | & | & \cdots & | \\ \mathbf{s}_1 & \mathbf{s}_2 & \cdots & \mathbf{s}_M \\ | & | & \cdots & | \end{bmatrix}$$

where the  $\{\mathbf{s}_k\}$  are codewords of dimension  $N$ . We usually assume  $M \geq N$  but that's written in stone. We also have

$$\begin{aligned} \mathbf{S}\mathbf{S}^\top &= \mathbf{S}\mathbf{S}^\top \\ \mathbf{R} &= \mathbf{S}\mathbf{S}^\top + \mathbf{W} \\ C_s &= \frac{1}{2} \log |\mathbf{R}| + \frac{1}{2} \log |\mathbf{W}| \end{aligned}$$

Greedy interference avoidance applied to codeword  $k$  replaces  $\mathbf{s}_k$  with  $\mathbf{x}$  where  $\mathbf{x}$  is a minimum eigenvalue eigenvector of  $\mathbf{R} - \mathbf{s}_k \mathbf{s}_k^\top$  and we derived a bunch of conditions for fixed

points and optimality in class and in the paper, especially in relation to a quantity called the TSC (or GSC), defined as

$$\text{TSC} = \text{Trace} \left[ \left( \mathbf{S}\mathbf{S}^\top \right)^2 \right]$$

$$\text{GSC} = \text{Trace} \left[ \mathbf{R}^2 \right]$$

- (a) We know that in white noise, greedy interference avoidance does not increase TSC. Show that greedy interference avoidance does not increase GSC in colored noise ( $\mathbf{W} \neq \alpha \mathbf{I}$ ).
- (b) Show that greedy interference avoidance does not decrease sum capacity.
- (c) Write an interference avoidance algorithm (in Matlab or C). You may assume white unit variance noise  $\mathbf{W} = \mathbf{I}$ . Starting from randomly chosen unit norm codewords, sketch (plot) the codeword set attained by interference avoidance for  $M = 3$  users in  $N = 2$  dimensions. Do the same for  $M = 8$  users in  $N = 2$  dimensions. Can you find examples of initial codeword sets (for  $M = 3$ ) which do not attain minimum TSC even after repeated application of your algorithm?

Plot the codeword covariance spectra in each case.

4. **Multi-Base Interference Avoidance** Now, suppose we have a TWO base system and two user groups with codeword matrices  $\mathbf{S}_1$  and  $\mathbf{S}_2$ . The users are arranged so that the covariance at base 1 is

$$\mathbf{R}_1 = \mathbf{S}_1\mathbf{S}_1^\top + g\mathbf{S}_2\mathbf{S}_2^\top + \mathbf{I}$$

and

$$\mathbf{R}_2 = \mathbf{S}_2\mathbf{S}_2^\top + g\mathbf{S}_1\mathbf{S}_1^\top + \mathbf{I}$$

at base 2 where  $g$  is a non-negative gain factor representing the interference level seen at base  $i$  from user set  $j$  (a symmetric interference system). Assume  $N = 3$  and  $M = 3$  for each set of users (for a total of six users over the two bases).

Write code which applies greedy interference avoidance at each base. For  $g = 0.1$  and  $g = 2$  first try sequential application (one base then the other). Then try interleaved application and answer the following questions.

- (a) One could define two TSC's for this problem:

$$\text{TSC}_1 = \text{Trace} \left[ \left( \mathbf{S}_1\mathbf{S}_1^\top \right)^2 \right]$$

and

$$\text{TSC}_2 = \text{Trace} \left[ \left( \mathbf{S}_2\mathbf{S}_2^\top \right)^2 \right]$$

Plot each TSC on the same graph as a function of algorithm step. When (if ever) do either or both of the TSC's converge? Is there a pattern of convergence as a function of the interference gain factor  $g$ ?

- (b) If the codewords converge, plot the converged spectra for both user sets. Do the codeword spectra overlap in signal space? If so, when and how much?

## 5. Potential Research Paper Topics COMING!